## **NORMALIZATION**

Normalization is the process of organizing data in a database to:

Reduce redundancy

Avoid update, insert, and delete anomalies

Maintain consistency

Improve efficiency

# 1. FIRST NORMAL FORM (1NF)

## **Definition**

A relation is in **1NF** if:

All attributes contain **atomic** (**indivisible**) values.

No multivalued or repeating attributes.

Each record is unique.

# **Example (NOT in 1NF)**

# **Employee Table**

# **EmpID EmpName Skills** PhoneNumbers

101 Ravi Java, Python 98765, 87654

102 Priya SQL 76543

#### **Problems:**

Skills contains multiple values  $\rightarrow$  multivalued

PhoneNumbers contains multiple values

Violates 1NF

## **Convert to 1NF**

Split multi-valued attributes into separate rows.

# **Employee Table (1NF)**

# **EmpID EmpName Skill PhoneNumber**

101 Ravi Java 98765

101 Ravi Python 87654

102 Priya SQL 76543

Now all values are atomic  $\rightarrow$  **1NF achieved** 

## 2. SECOND NORMAL FORM (2NF)

## **Definition**

A table is in **2NF** if:

It is in 1NF

## No partial dependency exists

A non-key attribute should NOT depend on part of a composite key.

(F) Applies only if a table has a **composite primary key**.

# **Example of Partial Dependency**

OrderDetails(OrderID, ProductID → Composite Key)

## **OrderID ProductID ProductName Price Qty**

Functional Dependencies:

 $\{OrderID, ProductID\} \rightarrow Qty$ 

ProductID → ProductName, Price (partial dependency)

## **Fixing 2NF Violation**

Split into two tables:

## 1. Product Table

## **ProductID ProductName Price**

## 2. OrderDetails

# **OrderID ProductID Qty**

Now:

No partial dependency

Achieved 2NF

# 3. THIRD NORMAL FORM (3NF)

## **Definition**

A relation is in **3NF** if:

It is in 2NF

No **transitive dependency** exists:

 $X \rightarrow Y$  and  $Y \rightarrow Z$ means  $X \rightarrow Z$  is transitive

# **Example of Transitive Dependency**

## **Student Table**

# RollNo Name DeptID DeptName

Functional Dependencies:

 $RollNo \rightarrow DeptID$ 

DeptID → DeptName

So: RollNo → DeptName (transitive)

## Convert to 3NF

Split the table:

## 1. Student

| RollNo | Name | DeptID |

## 2. Department

| DeptID | DeptName |

Now:

No transitive dependency

Achieved 3NF

# 4. BOYCE-CODD NORMAL FORM (BCNF)

## **Definition**

A relation is in **BCNF** if:

For every functional dependency  $X \rightarrow Y$ ,

X must be a superkey.

More strict than 3NF.

# **BCNF** Example

Lecture Hall Allocation Table:

## **Course Instructor Room**

Assume:

Instructor → Room

Room → Instructor

(Course, Instructor)  $\rightarrow$  Room

## **Problem:**

Instructor is NOT a key

Room is NOT a key

Both determine other values → violates BCNF

# **Convert to BCNF**

Split:

## 1. Instructor-Room

| Instructor | Room |

## 2. Course-Instructor

| Course | Instructor |

Now FDs:

Instructor  $\rightarrow$  Room (Instructor is key)

Course → Instructor (Course is key)

BCNF achieved.

# **5. FOURTH NORMAL FORM (4NF)**

## **Definition**

A table is in 4NF if:

It is in BCNF

It has no non-trivial multivalued dependencies (MVDs)

# **Example of MVD**

A professor may teach multiple subjects

A professor may have multiple offices

# **Prof Subject** Office

A DBMS Block-1

A DBMS Block-2

A Networks Block-1

A Networks Block-2

Here:

 $\operatorname{Prof} \longrightarrow \operatorname{Subject}$ 

 $Prof \rightarrow \rightarrow Office$ 

Two independent multi-valued attributes → violates 4NF

#### Convert to 4NF

## 1. Prof-Subject

| Prof | Subject |

## 2. Prof-Office

| Prof | Office |

Now:

No multi-valued dependency in a single table

Achieved 4NF

# 6. FIFTH NORMAL FORM (5NF)

#### **Definition**

A relation is in **5NF** if:

It is in 4NF

It cannot be decomposed further without losing information

Deals with **join dependencies** 

Used when:

Data comes from three or more independent entities

## **Example**

Supplier-Product-Region table:

# **Supplier Product Region**

Assume:

A supplier can supply multiple products

A product can be supplied in multiple regions

Supplier works in multiple regions

These combine to create many rows.

To achieve 5NF, split into:

- 1. Supplier-Product
- 2. Product-Region
- 3. Supplier-Region

Now joins reconstruct the original table without anomalies.

# CASE STUDY 1: Normalizing a RETAIL INVENTORY DATABASE UNNORMALIZED RetailInventory Table

# ItemID ItemName Supplier SupplierPhone StoreLocations PurchasePrices

101 Pen ABC Co 98765 Chennai, Pune 10, 9.5

## **Problems**

StoreLocations is multi-valued

PurchasePrices is multi-valued

Supplier details repeated → redundancy

## **Step 1: Convert to 1NF**

Break multi-valued attributes:

# ItemID ItemName Supplier SupplierPhone StoreLocation PurchasePrice

101	Pen	ABC Co 98765	Chennai	10
101	Pen	ABC Co 98765	Pune	9.5

# **Step 2: Remove Partial Dependencies (2NF)**

Assume Key = {ItemID, StoreLocation}

But Supplier depends only on ItemID  $\rightarrow$  partial dependency.

Split:

#### Item

| ItemID | ItemName | Supplier |

# **ItemLocation**

| ItemID | StoreLocation | PurchasePrice |

# **Supplier**

| Supplier | SupplierPhone |

## **Step 3: Remove Transitive Dependencies (3NF)**

Supplier → SupplierPhone is fine No transitive dependencies remain.

#### Final BCNF / 4NF / 5NF

No multivalued dependencies.

Final database:

- ✓ Item(ItemID, ItemName, SupplierID)
- **✓** Supplier(SupplierID, SupplierName, SupplierPhone)
- **✓** ItemLocation(ItemID, StoreLocation, PurchasePrice)

# **CASE STUDY 2: Normalizing a LIBRARY DATABASE**

# **UNNORMALIZED Library Table**

**BookID Title Authors Member BorrowDate** 

B101 DBMS Navathe, Korth Suresh 2023-01-01

#### **Problems**

Authors is multi-valued

Book borrowed by a member repeatedly → redundancy

## 1NF

Split authors:

# **BookID Title Author Member BorrowDate**

# 2NF & 3NF

Dependencies:

BookID  $\rightarrow$  Title

BookID  $\rightarrow$  Author (M:N relationship  $\rightarrow$  separate table)

Member → BorrowDate depends on (BookID, Member)

# **Final 3NF Schema**

- **✓** Book(BookID, Title)
- **✓** Author(AuthorID, AuthorName)
- **✓** BookAuthor(BookID, AuthorID)
- **✓** Member(MemberID, MemberName)
- **✓** Borrow(BookID, MemberID, BorrowDate)

# FINAL SUMMARY TABLE OF NORMAL FORMS

Normal Form	Removes	Violation Example
1NF	Multivalued attributes, repeating groups	Skills = {Java, Python}
2NF	Partial dependency	ProductID → ProductName
3NF	Transitive dependency	$DeptID \rightarrow DeptName$
BCNF	Determinant not a key	$Instructor \rightarrow Room$
4NF	Multivalued dependency	$\operatorname{Prof} \longrightarrow \operatorname{Subject}$ , Office
5NF	Join dependency	Supplier-Product-Region