10. Operand addressing

An addressing mode is a method of specifying the data source or destination in an instruction. There are 5 types of addressing modes is supported by 8051.

- 1. Register
- 2. Immediate
- 3. Direct (memory related)
- **4.** Register Indirect (memory related)
- 5. Index register addressing

Register addressing mode

This addressing mode involves the use of registers to hold the data to be manipulated.

Examples:

MOV A, R0; Copy the contents of R0 int A

ADD A, R7; Add the contents of R7 to contents of A and the result is stored in A

Immediate addressing mode

In this addressing mode immediate data is specified in instruction as a source operand.

Examples:

MOV B, #40H ; load 40H into B register

MOV DPTR, #2000H ; load 2000H into DPTR

Direct addressing mode

As we know the on-chip RAM of 8051 is 128 byte, it can be accessed through memory address from 00H to FF H. The allocations of 128 bytes are as follows.

- 1. RAM location 00H-1FH are assigned to register banks and stack
- 2. RAM location 20H-2FH is set aside as bit-addressable space to save single bit data.
- 3. RAM location 30H-7F is available as place to save bite-sized data.

Although the entire 128 bytes of RAM can be accessed through direct addressing mode, it is most often used to access RAM location 30H-7FH. This is due to fact that register banks are accessed through their names.

Examples:

MOV R4, 70H; move the contents of RAM location 70H to R4.

MOV 56H, A ; save the content of A in RAM location 56H

PUSH 05; push R5 onto the stack

Register indirect addressing mode

In this mode the address (of 8bits) is indirectly specified in the instruction by the contents of pointer. This addressing mode so called because the source operand is from the address specified indirectly by another register in the instruction. The limitation is that only R0 and R1 register can be used in 8051 for indirect addressing. SFRs are directly accessible.

Examples

MOV R1, #55H ; load pointer R1=55H

MOV A, @R1 ; the content of pointer is transferred to A

Index registers addressing

Suppose we need to access external data RAM and external code space of on-chip ROM 16 bit address must be required. In this case we have to use DPTR. This mode is widely used in accessing data elements of look-up table entries in the program ROM space of 8051.

Examples;

MOV DPTR, #0200H ; load DPTR with 0200

CLR A ; clear accumulator

MOVC A,@A+DPTR ; Move the content 0200 location into A

11. Instruction set

The instruction set of 8051 can be classified into following group.

- 1. Data Transfer Instructions
- 2. Arithmetic Instructions
- 3. Logic Instructions
- 4. Boolean Variable manipulation Instructions
- 5. Program flow control (Processor and Machine control) Instructions
- 6. Interrupt flow Control instruction

12.1 Data Transfer Instruction

Three types of the data transfer can be done by move instruction. First type is transfer within the internal RAM and SFRs, second type is transfer using code memory area (CODE) and the third is using the external data memory X-DATA).

MOV instruction

A MOV instruction means move (copy) the bits from one source to a destination.

Table 4.4 MOV instructions within the registers, internal RAM and SFRs in 8051

Instruction	Action	Addressing	Length	cycles
(Mnemonic)			in bytes	
MOV A, Rn	Move Rn into A	Register	1	1
MOV Rn, A	Move into Rn from A	Register	1	1
MOV A, #data	Move immediate 8-bit data into A	Immediate	2	1
MOV Rn, #data	Move into Rn the data.	immediate	2	1
MOV A, direct	Move byte at the direct address into A	Direct	2	1
MOV Rn, direct	Move from direct address into Rn	Direct	2	2
MOV direct, A	Move byte to the direct address form A	Direct	2	1
MOV direct, Rn	Move a byte to the direct address from Rn	Direct	2	2
M OV direct, direct	Move byte to the direct address from the direct address	Direct	3	2
MOV direct, #data	Move immediate data byte to the direct address	Immediate	3	2
MOV a,@Ri	Move into A the byte from the address pointed by Ri	Indirect	2	2
MOV @Ri, A	Move A into address pointed by Ri	Indirect	1	1
MOV direct, @Ri	Move into direct address from address pointed by Ri	indirect	1	1
MOV @Ri, direct	Move from the direct address to the address poined by ri	Indirect	2	2
MOV @Ri, #data	Move data ino address pointed by Ri	immediate	2	2
MOV DPTR, data16	Mov e16 bit dat	immediate	3	2

MOVC-type Instruction

It moves the 8-bit code from one source at the program memory (internal and external) to the register A destination.

Table 4.5 MOVC Instructions for transfer from the program memory area address code or constant to accumulator in 8051

Instruction	Action	Addressing	Length in bytes	Cycles
MOVC A, @A+DPTR	Moves the code or constant into A the byte	Indirect	1	2
	from the program memory address pointed			

	by hypothetical addition of DPTR with the A itself.			
MOVC A, @A+PC	Move the code or constant into A the byte from the program memory address pointed by hypothetical addition of PC with the A itself	Indirect	1	2

MOX-type Instructions

A MOVX instruction means move (copy) the 8-bit data into A and from A using the external data memory address using DPTR or Ri as the pointer

Table 4.6 MOVX instruction

Instruction	Action	Addressing	Length in bytes	Cycles
MOVX A, @DPTR	Move the external data byte (X-DATA) into A from the data memory address pointed by DPTR	Indirect	1	2
MOVX @DPTR,A	Move into the external data memory from A to the address pointed by DPTR	Indirect	1	2
MOVX A,@Ri	Move the external data byte into a from the memory address pointed by Ri	Indirect	1	2
MOVX @Ri, A	Move into the external data memory from A to the memory address pointed by Ri	Indirect	1	2

Table 4.7 PUSH and POP instructions for using the Stack Area employing SP

Instruction	Action	Addressing	Length in	Cycles
			bytes	
PUSH direct	Move byte from a direct internal RAM or SFR into the stack after first incrementing the stack pointer by 1	Direct	2	2
POP direct	Move byte to a direct internal RAM or SFR into the stack and then decrement the stack pointer by 1.	Direct	2	2

XCH-type instructions

An XCH instruction is for exchanging the A register with a source using the register (direct or indirect addresing) mode.

Table 4.8 XCH and XCHD instruction

Instruction	Action	Addressing	Length in bytes	cycles
XCH A@Ri	Exchange byte at A with the address pointed by Ri	Indirect	1	2
XCH A,Rn	Exchange byte at A with the register Rn	Register	1	2
XCH A, direct	Exchange byte at A with the byte at a direct address.	Direct	1	1
XCHD A,@Ri	Exchange lower hex-digits of the bytes at A with the address pointed by Ri	Indirect	1	2

12.2 Arithmetic Instruction

These instructions include 8 bit addition, subtraction, increment, decrement, multiply and division instruction.

Table 4.9 Arithmetic ADD, SUB, MUL, DIV, INC and DEC instruction s in 8051

Instruction	Action	Addressing	Flags	Length	Cycles
			affected	(bytes)	
ADD A,Rn	Add Rn into A	Register	C,AC,OV	1	1
ADD A, direct	Add the byte at the direct address into A	Direct	C,AC,OV	2	1
ADD A, @Ri	Add the byte from the address pointed by the Ri into A	Indirect	C,AC,OV	1	1
ADD A, #data	Add immediate data byte to the A	Immediate	C,AC,OV	2	1
ADDC A, Rn	Add CF(carry) bit and Rn into A	Register	C,AC,OV	1	1
ADDC A, direct	Add CF bit and byte at the direct address ito A	Direct	C,AC,OV	2	1
ADDC A @Ri	Add CF bit and the byte from the address pointed by the Ri	Indirect	C,AC,OV	1	1
ADDC A, #data	Add CF bit and immediate data byte to the A	Immediate	C,AC,OV	2	1
SBBB A,Rn	Subtract borrow at CF bit and Rn into A	Rgister	C,AC,OV	1	1
SBBB A, direct	Subtract borrow at CF bit and byte at the direct address into A	Direct	C,AC,OV	2	1
SBBB A, @Ri	Subtract borrow at C bit and byte at the byte from the address pointed	Indirect	C,AC,OV	1	1

	by the Ri into A				
SBBB A, #data	Subtract borrow at CF bit and	Immediate	C,AC,OV	2	1
	immediate data byte into A				
INC A	Increment	Register	None	1	1
INC Rn	Increment Rn	Register	None	1	1
INC direct	Increment byte at the direct address	Direct	None	2	1
INC @Ri	Increment the byte at the address	Indirect	None	1	1
	pointed by Ri				
DEC A	Decrement A	Register	None	1	1
DEC Rn	Decrement Rn	Register	None	1	1
DEC direct	Decrement byte at the direct	Direct	None	2	1
	address				
DEC @Ri	Decrement the byte at the address	Indirect	None	1	1
	pointed by the Ri				
MUL AB	Multiply A and B Result MSB in B	Register	OV	1	4
	and LSB in A				
DIV AB	Divide A (Numerator) and B(Register	OV	1	4
	denominator) Remainder in B				
	Quotient in A				
DAA	Decimal adjust accumulator	Register	С	1	1

12.3 Logical Instruction

Table gives features of 8-bit AND, OR and XOR instruction. These instructions have 4 addressing modes such as register, immediate, direct and indirect.

Table 4.10 ANL, ORL XRL instruction

Instruction	Action	Addressing	Length in bytes	Cycles
ANL A, Rn	AND Rn into A	Register	1	1
ANL A, direct	AND byte at the direct address into A	Direct	2	1
ANL A, @Ri	AND into the byte from the address pointed by the Ri	Indirect	1	1
ANL A, #data	AND immediate data byte into A	immediate	2	1
ANL direct, A	AND A into byte at the direct address	Direct	2	1
ANL direct, #data	AND immediate byte into byte at the direct address	Direct	3	2
ORL A, Rn	OR Rn into A	Register	1	1
ORL A, direct	OR byte at the direct address into A	Direct	2	1
ORL A, @Ri	OR into the byte from the address pointed by Ri	Indirect	1	1
ORL A, #data	OR immediate data byte to the A	immediate	2	1
ORL direct, A	OR A into byte at the direct address	Direct	2	1
ORL direct,#data	OR immediate byte into byte at the	Direct	3	2

	direct address			
XRL A, Rn	XOR Rn into A	Register	1	1
XRL A, direct	XOR byte at the direct address	Direct	2	1
	into A			
XRL A, @Ri	XOR the byte at the address	Indirect	1	1
	pointed by Ri into A			
XRL A, #data	XOR immediate data byte to the A	immediate	2	1
XRL direct, A	XOR A into byte at the direct	Direct	2	1
	address			
XRL direct, #data	XOR immediate byte into byte at	Direct	3	2
	the direct address			

12.4 Boolean Variable manipulation Instructions

These are also called as Boolean processing instruction.

Table 4.11 MOV, CLR, CPL, SETB, ANL, and ORL Boolean Processing Instruction

Instruction	Action	Addressing	Length (bytes)	Cycles
MOV C, bit	Move bit into CF	Direct bit addressing	2	1
MOV bit, C	Move CF into the bit	Direct bit addressing	2	2
CLR C	Clear CF	PSW Register CF bit	1	1
		addressing		
CLR bit	Clear bit	Direct bit addressing	2	1
CPL C	Complement CF	PSW Register CF bit	1	1
		addressing		
CPL bit	Complement bit	Direct bit addressing	2	1
SETB C	Set CF=1	PSW Register CF bit	1	1
		addressing		
SETB bit	Set bit =1	Direct bit addressing	2	1
ANL C,bit	AND between CF and bit, place the result in CF	Direct bit addressing	2	2
ANL C, bit	AND between CF and , place the result in C	Direct bit addressing	2	2
ORL C,bit	OR between CF and bit, place the result in C	Direct bit addressing	2	2
ORL C, bit	OR between CF and \overline{bit} , place the result in C	Direct bit addressing	2	2

12.5 Control Transfer Instruction

In the main program other sub programs may be called to perform a particular task. When a sub program is called the processor will jump to a new address where this program is available and

it has to accomplish program flow control transfer with help of JUMP and CALL instruction when some condition met.

Table 4.12 Delay-Cycle (NOP) instruction (No operation)

Instruction	Action	Addressing	Length in	Cycles
			bytes	
NOP	No operation, PC gets the address of		1	1
	next instruction on incrementing at			
	NOP.			

Long, Absolute and Short Jump

8051 has three jump instructions: Long- it jumps to 16-bit address, Absolute- it jumps within 2 K bytes and Short- it jumps to address within 128 bytes above or below the present address.

Table 4.13 Long, absolute and short jump instructions

Instruction	Action	Addressing	Length in bytes	Cycles
LJMP addr16	Jump to the next address given by two bytes in the instruction	Direct 16 bit address	•	2
AJMP addr11	Jump to the next address	Direct 11-bit address	2	2
SJMP rel	Jump in the range between -128 and +127 from the address of next instruction	Direct 8-bit	2	2
JMP @A+DPTR	Jump in the next address given by addition of 8-bits of A with 16-bits of DPTR	Indirect 16-bit relative addreess		

Table 4.14 Conditional Short Relative Jumps

Instruction	Action	Addressing	Length in bytes	Cycles
JNZ rel	Jump to a relative address if a is not zero	Relative(offset)	2	2
JZ rel	Jump to a relative address if A is zero	Relative(offset)	2	2
JNC rel	Jump to a relative address if CF is not 1	Relative(offset)	2	2
JC rel	Jump to a relative address if CF=1	Relative(offset)	2	2
JB bit, rel	Jump to a relative address if addressed bit 1 (bit not set)	Relative(offset)	2	2
JNB bit,rel	Jump to a relative address if addressed bit 0 (bit not set)	Relative(offset)	2	2
JBC bit, rel	Jump to a relative address if addressed bit 1(bit set) and reset	Relative(offset)	2	2

(1 CF 0)		
Loorry (moko ('E-II)		
carry (make cr-0)		

Decrement and Conditional jump on Zero

Table 4.15 Instruction for decrement and then jump in program-loops in 8051

Instruction	Action	Addressing	Length in bytes	Cycles
DJNZ Rn, Rel	Decrement Rn and jump if Rn is still not zero.	Relative (offset)	2	2
DJNZ direct, Rel	Decrement byte at the direct and jump if byte is still not zero	Relative (offset)	2	2

Jump after comparison

Table 4.16 Compare then conditional jump after comparison

Instruction	Action	Addressing	Flag	Length	Cycles
			affected	in bytes	
CJNE A, #data, rel	Compare A and	Relative	C	3	2
	immediate data and	(offset)			
	jump if both are not				
	equal.				
CJNE Rn, #data, rel	Compare Rn and	Relative	С	3	2
	immediate data and	(offset			
	jump if both are not				
	equal.				
CJNE A, direct, rel	Compare the bytes at	Relative	C	3	2
	A and direct and	(offset			
	jump if both are not				
	equal				
CJNE @Ri, #data, rel	Compare byte from	Relative	C	3	2
	the address pointed	(offset)			
	by Ri and immediate				
	data and jump if				
	both are not equal				

Call to a Routine

`Table 4.17 Long, absolute call and return instruction

Instruction	Action	Addressing	Length in bytes	Cycles
LCALL addr16	Call to the next address given by two	Direct 16-	3	2
	bytes in the instruction	bit address		
ACALL addr11	Call the next address given by 11 bits in	Direct 11	2	2

	the instruction.	bit address		
RET	Return to PC the saved PCL and PCH	Stack	1	2
	from the stack.	address		

12.6 Interrupt Control Flow (RETI instruction)

Table 4.18 RETI instruction

Instruction	Action	Addressing	Length In bytes	cycles
RETI	Return into PC the	Stack adddress	1	2
	saved PCL and			

12. Programming

While the CPU can work only in binary, it can do so at a very high speed, however, it is quite tedious and slow for humans to deal with 0s and 1s in order to program the computer. A program that consists of 0s and 1s is called machine language. In the early days of the computer programmers coded programs in machine language. Although the hexadecimal system was used as a more efficient way to represent binary numbers, the process of working in machine code was still cumbersome for humans. Eventually, assembly language were developed which provided **mnemonics** for the machine code instructions. Plus other features which made programming faster and less prone to error. Assembly language is referred to as low level language because it deals directly with internal structure of CPU. Programmer needs assembler to convert the assembly language to machine language for execution purpose. Assembly language consists mnemonics optionally followed by one or two operands.

Programs

P1. Write an ALP (Assembly Language Program) to find the sum of values and store the result in A (lower byte and in R7 (higher byte). Assume that RAM locations 40-44 have the following values.

Solution:

MOV R0, #40H ; load pointer

MOV R2, #05H ; load counter

CLR A ; A=0

MOV R7, A ; clear R7

AGAIN: ADD A, @RO ; add the byte pointer

JNC NEXT; if CY=0 it can jump to NEXT label

INC R7; increment counter

NEXT: INC R0 ; increment pointer

DJNZ R2,AGAIN ; repeat until R2is zero

HERE: SJMP HERE

P2. Assume that 5 BCD data items are stored in RAM locations starting a 40H as shown below. Write an ALP to find the sum of all numbers. The result must be in BCD.

Solution:

MOV R0, #40H ; load pointer

MOV R2, #05H ; load counter

CLR A ; A=0

MOV R7, A ; clear R7

AGAIN: ADD A, @RO; add the byte pointer

DA A

JNC NEXT ; if CY=0 it can jump to NEXT label

INC R7; increment counter

NEXT: INC R0 ; increment pointer

DJNZ R2,AGAIN ; repeat until R2is zero

HERE: SJMP HERE

P3. Write an ALP to get hex data in the range of 00-FFH from port 1 and convert it to decimal. Save the digits in R7, R6 and R5, where the least significant digit in R7.

MOV A, #0FFH

MOV P1, A ; make an P1 an input port

MOV A1, P1; read data from P1

MOV B, #0AH ; move 0AH to register b

DIV AB; divide by the contents of A by B

MOV R7, B ; Save lower digit in R7 register

MOV B, #0AH ;

DIV AB ;

MOV R6, B; save the next digit

MOV R5, A ; save the last digit

HERE: SJMP HERE

P4. Read and test P1 to see whether it has the value 45H. if it does send 99H to P2; otherwise, it stays cleared.

Solution:

MOV P2, 00H ; clear P2

MOV P1, #0FFH ; make P1 an input port

MOV R3, #45H ; R3=45H

MOV A, P1 ; read P1

XRL A, R3;

JNZ EXIT

MOV P2, #99H

EXIT:

P5. Find the 2's complement of the value 78 H

Solution:

MOV A, #78H ; A=85H

CPL A ; make 1's complement a

ADD A, #01H ; make 2's complemt

HERE: SJMP HERE

P6. Write an ALP to determine if register A contains the value 99H, if so, make R1=FFH otherwise make R1=0.

Solution:

MOV R1, #00H ; clear R1

CJNE A, #99H, NEXT ; if A is not equal 99H then jump

MOV R1, #0FFH ; make R1=FFH

NEXT

P7. Assume that P1 is an input port connected to a temperature sensor. Write an ALP to read the temperature and test it for the value 75. According to the rest result, place the temperature value into the registers indicated by the following.

If T=75 then A=75

If T < 75 then R1 = T

If T>75 then R2=T

Solution:

MOV P1, # 0FFH; make P1 an input port

MOV A, P1 ; read P1 port, temperature

CJNE A, #75, OVER ; jump if A is not equal 75

SJMP EXIT

OVER: JNC NEXT ; if CY=0, then A>75

MOV R1, A ; if CY=1, A<75

SJMP EXIT ; Exit

NEXT: MOV R2, A

EXIT

P8. Write an ALP that finds the number of 1s in a given byte 97H.

Solution:

MOV R1, #00H ; clear R1

MOV R7, 08H ; Counter=08

MOV A, 97H

AGAIN: RLC A ; rotate through CY once

JNC NEXT ; check for CY

INC R1 ; if CY=1 then increment R1

NEXT: DJNZ R7, AGAIN ; go through 8times

HERE: SJMP HERE

P9. Assume that register a has packed BCD 29H, write an ALP to convert packed BCD to ASCII numbers and place them in R2 and R6.

Solution:

MOV A, #29H ; A=29H, packed BCD

MOV R2, A ; keep a copy of BCD data in R2

ANL A, #0FH ; mask the upper nibble (A=09)

ORL A, #30H ; make it an ASCII, A=39H

MOV A, R6 ; save in R6

MOV A, R2 ; A=29H

ANL A, #0F0H; mask the lower nibble

RR A ; rotae right

RR A ; rotae right

RR A ; rotate right

RR A ; rotate right

ORL A, #30 H ; A=32H

MOV R2, A ; save the ASCII character in R2

HERE: SJMP HERE

P10. Write an ALP to create a square wave of 50% duty cycle on bit 0 of port 1.

Solution:

HERE: SETB P1.0 ; set to high bit 0 of port 1

LCALL DELAY ; call the delay subroutine

CLR P1.0 ; p1.0=0

LCALL DELAY

SJMP HERE

P11. Assume that the bit P2.2 is used to control the outdoor light and bit P2.5 to control the light inside the building. Write an ALP to turn on outside light and to turn the inside one.

Solution:

SETB C ; CY=1
ORL C, P2.2 ; CY=P2.2

MOV P2.2, C ; turn it "on" if not already "on"

CLR C ; CY=0

ANL C, P2.5 ; CY=P2.5 ANDed with CY

MOV P2.5, C ; turn it off if not already off.